

## Chapter 6

# Recursion as a Problem-Solving Technique

# Backtracking

- Backtracking
  - A strategy for guessing at a solution and backing up when an impasse is reached
- Recursion and backtracking can be combined to solve problems

# The Eight Queens Problem

- Problem
  - Place eight queens on the chessboard so that no queen can attack any other queen
- Strategy: guess at a solution
  - There are 4,426,165,368 ways to arrange 8 queens on a chessboard of 64 squares

# The Eight Queens Problem

- An observation that eliminates many arrangements from consideration
  - No queen can reside in a row or a column that contains another queen
    - Now: only 40,320 arrangements of queens to be checked for attacks along diagonals

# The Eight Queens Problem

- Providing organization for the guessing strategy
  - Place queens one column at a time
  - If you reach an impasse, backtrack to the previous column

# The Eight Queens Problem

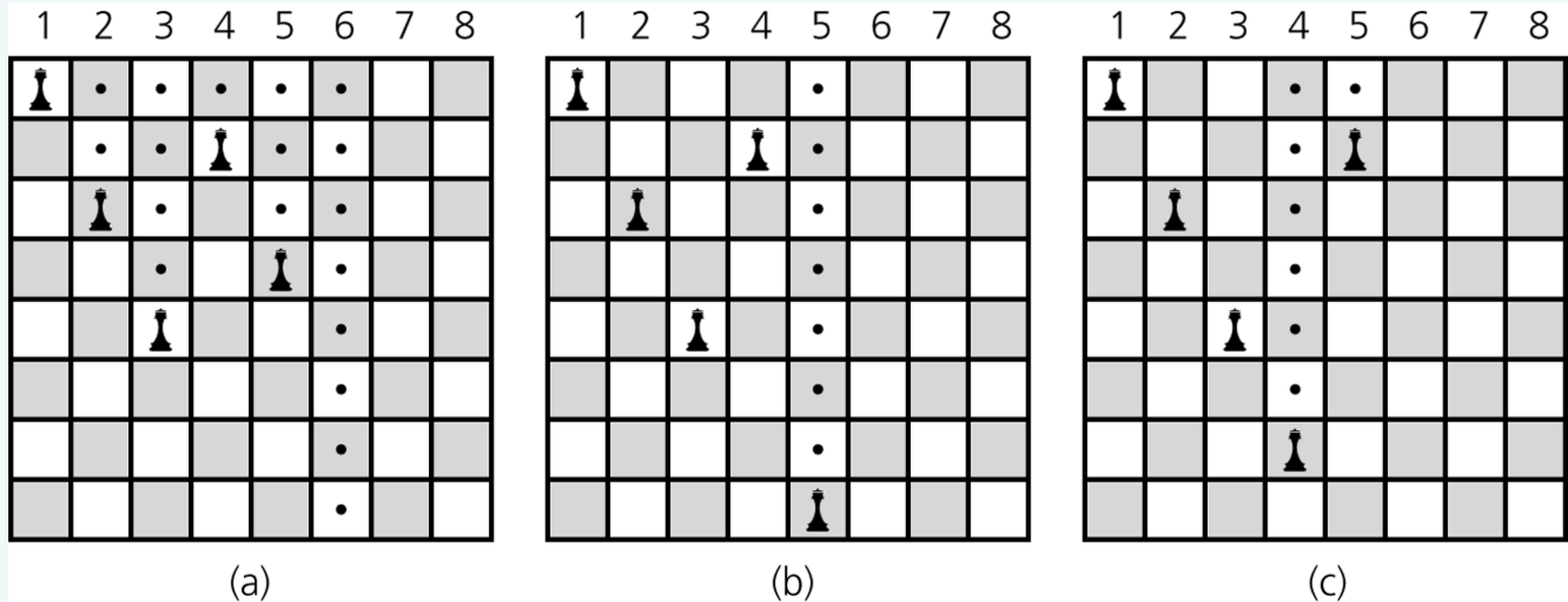


Figure 6-1

a) Five queens that cannot attack each other, but that can attack all of column 6; b) backtracking to column 5 to try another square for the queen; c) backtracking to column 4 to try another square for the queen and then considering column 5 again

# The Eight Queens Problem

- A recursive algorithm that places a queen in a column
  - Base case
    - If there are no more columns to consider
      - You are finished
  - Recursive step
    - If you successfully place a queen in the current column
      - Consider the next column
    - If you cannot place a queen in the current column
      - You need to backtrack

# The Eight Queens Problem

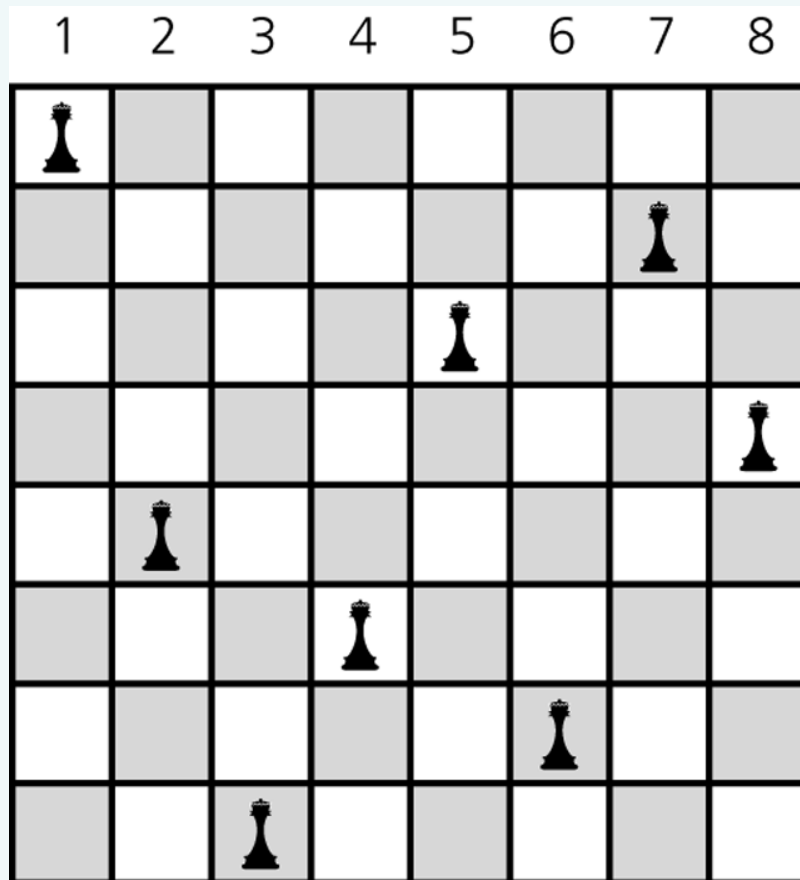


Figure 6-2

A solution to the Eight Queens problem



# Defining Languages

- A language
  - A set of strings of symbols
  - Examples: English, Java
  - If a Java program is one long string of characters, the language JavaPrograms is defined as
$$\text{JavaPrograms} = \{\text{strings } w : w \text{ is a syntactically correct Java program}\}$$

# Defining Languages

- A language does not have to be a programming or a communication language

- Example

- The set of algebraic expressions

- $$\text{AlgebraicExpressions} = \{w : w \text{ is an algebraic expression}\}$$

# Defining Languages

- Grammar
  - States the rules for forming the strings in a language
- Benefit of recursive grammars
  - Ease of writing a recognition algorithm for the language
    - A recognition algorithm determines whether a given string is in the language

# The Basics of Grammars

- Symbols used in grammars
  - $x \mid y$  means  $x$  or  $y$
  - $xy$  means  $x$  followed by  $y$   
(In  $x \bullet y$ , the symbol  $\bullet$  means concatenate, or append)
  - $\langle \text{word} \rangle$  means any instance of word that the definition defines

# The Basics of Grammars

- Java identifiers
  - A Java identifier begins with a letter and is followed by zero or more letters and digits

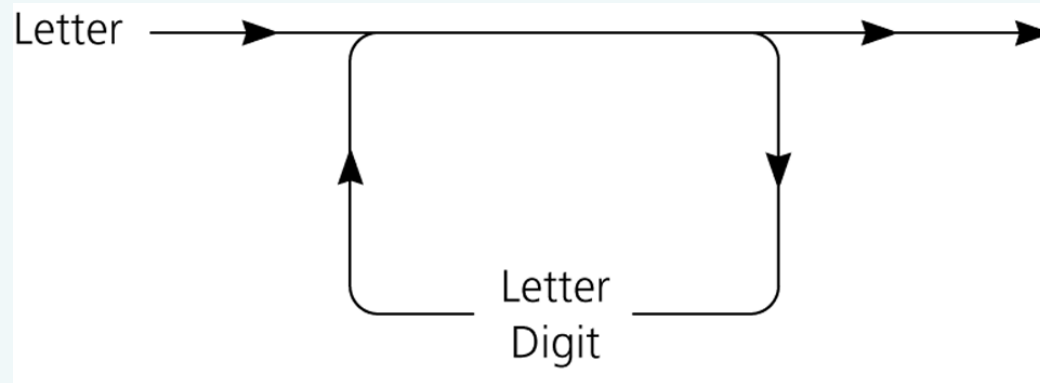


Figure 6-3

A syntax diagram for Java identifiers

# The Basics of Grammars

- Java identifiers

- Language

$\text{JavaIds} = \{w : w \text{ is a legal Java identifier}\}$

- Grammar

$\langle \text{identifier} \rangle = \langle \text{letter} \rangle \mid \langle \text{identifier} \rangle \langle \text{letter} \rangle \mid \langle \text{identifier} \rangle \langle \text{digit} \rangle$

$\langle \text{letter} \rangle = a \mid b \mid \dots \mid z \mid A \mid B \mid \dots \mid Z \mid \_ \mid \$$

$\langle \text{digit} \rangle = 0 \mid 1 \mid \dots \mid 9$

# The Basics of Grammars

- Recognition algorithm

```
isId(w)
  if (w is of length 1) {
    if (w is a letter) {
      return true
    }
    else {
      return false
    }
  }
  else if (the last character of w is a letter or a digit) {
    return isId(w minus its last character)
  }
  else {
    return false
  }
```

# Two Simple Languages: Palindromes

- A string that reads the same from left to right as it does from right to left
- Examples: radar, deed
- Language

Palindromes =  $\{w : w \text{ reads the same left to right as right to left}\}$



# Palindromes

- Grammar

$\langle \text{pal} \rangle = \text{empty string} \mid \langle \text{ch} \rangle \mid a \langle \text{pal} \rangle a \mid b \langle \text{pal} \rangle b \mid \dots$   
 $\mid Z \langle \text{pal} \rangle Z$

$\langle \text{ch} \rangle = a \mid b \mid \dots \mid z \mid A \mid B \mid \dots \mid Z$

# Palindromes

- Recognition algorithm

isPal(w)

```
if (w is the empty string or w is of length 1) {  
    return true  
}  
else if (w's first and last characters are the  
        same letter ) {  
    return isPal(w minus its first and last  
                characters)  
}  
else {  
    return false  
}
```

# Strings of the form $A^nB^n$

- $A^nB^n$ 
  - The string that consists of  $n$  consecutive  $A$ 's followed by  $n$  consecutive  $B$ 's
- Language
  - $L = \{w : w \text{ is of the form } A^nB^n \text{ for some } n \geq 0\}$
- Grammar
  - $\langle \text{legal-word} \rangle = \text{empty string} \mid A \langle \text{legal-word} \rangle B$

# Strings of the form $A^nB^n$

- Recognition algorithm

```
isAnBn(w)
    if (the length of w is zero) {
        return true
    }
    else if (w begins with the character A and ends
              with the character B) {
        return isAnBn(w minus its first and last
                       characters)
    }
    else {
        return false
    }
```

# Algebraic Expressions

- Three languages for algebraic expressions
  - Infix expressions
    - An operator appears between its operands
    - Example:  $a + b$
  - Prefix expressions
    - An operator appears before its operands
    - Example:  $+ a b$
  - Postfix expressions
    - An operator appears after its operands
    - Example:  $a b +$

# Algebraic Expressions

- To convert a fully parenthesized infix expression to a prefix form
  - Move each operator to the position marked by its corresponding open parenthesis
  - Remove the parentheses
  - Example
    - Infix expression:  $((a + b) * c$
    - Prefix expression:  $* + a b c$

# Algebraic Expressions

- To convert a fully parenthesized infix expression to a postfix form
  - Move each operator to the position marked by its corresponding closing parenthesis
  - Remove the parentheses
  - Example
    - Infix form:  $((a + b) * c)$
    - Postfix form:  $a b + c *$

# Algebraic Expressions

- Prefix and postfix expressions
  - Never need
    - Precedence rules
    - Association rules
    - Parentheses
  - Have
    - Simple grammar expressions
    - Straightforward recognition and evaluation algorithms



# Prefix Expressions

- Grammar

$\langle \text{prefix} \rangle = \langle \text{identifier} \rangle \mid \langle \text{operator} \rangle \langle \text{prefix} \rangle \langle \text{prefix} \rangle$

$\langle \text{operator} \rangle = + \mid - \mid * \mid /$

$\langle \text{identifier} \rangle = a \mid b \mid \dots \mid z$

- A recognition algorithm

```
isPre()
```

```
    size = length of expression strExp
```

```
    lastChar = endPre(0, size - 1)
```

```
    if (lastChar >= 0 and lastChar == size-1 {
```

```
        return true
```

```
    }
```

```
    else {
```

```
        return false
```

```
    }
```

# Prefix Expressions

- An algorithm that evaluates a prefix expression

```
evaluatePrefix(strExp)
    ch = first character of expression strExp
    Delete first character from strExp
    if (ch is an identifier) {
        return value of the identifier
    }
    else if (ch is an operator named op) {
        operand1 = evaluatePrefix(strExp)
        operand2 = evaluatePrefix(strExp)
        return operand1 op operand2
    }
```

# Postfix Expressions

- Grammar

$\langle \text{postfix} \rangle = \langle \text{identifier} \rangle \mid \langle \text{postfix} \rangle \langle \text{postfix} \rangle \langle \text{operator} \rangle$

$\langle \text{operator} \rangle = + \mid - \mid * \mid /$

$\langle \text{identifier} \rangle = a \mid b \mid \dots \mid z$

- At high-level, an algorithm that converts a prefix expression to postfix form

```
if (exp is a single letter) {
```

```
    return exp
```

```
}
```

```
else {
```

```
    return postfix(prefix1) + postfix(prefix2) +  
        operator
```

```
}
```

# Postfix Expressions

- A recursive algorithm that converts a prefix expression to postfix form

```
convert(pre)
    ch = first character of pre
    Delete first character of pre
    if (ch is a lowercase letter) {
        return ch as a string
    }
    else {
        postfix1 = convert(pre)
        postfix2 = convert(pre)
        return postfix1 + postfix2 + ch
    }
```

# Fully Parenthesized Expressions

- To avoid ambiguity, infix notation normally requires
  - Precedence rules
  - Rules for association
  - Parentheses
- Fully parenthesized expressions do not require
  - Precedence rules
  - Rules for association

# Fully Parenthesized Expressions

- Fully parenthesized expressions
  - A simple grammar
$$\langle \text{infix} \rangle = \langle \text{identifier} \rangle \mid (\langle \text{infix} \rangle \langle \text{operator} \rangle \langle \text{infix} \rangle)$$
$$\langle \text{operator} \rangle = + \mid - \mid * \mid /$$
$$\langle \text{identifier} \rangle = a \mid b \mid \dots \mid z$$
  - Inconvenient for programmers

# The Relationship Between Recursion and Mathematical Induction

- A strong relationship exists between recursion and mathematical induction
- Induction can be used to
  - Prove properties about recursive algorithms
  - Prove that a recursive algorithm performs a certain amount of work

# The Correctness of the Recursive Factorial Method

- Pseudocode for a recursive method that computes the factorial of a nonnegative integer  $n$

```
fact(n)
  if (n is 0) {
    return 1
  }
  else {
    return n * fact(n - 1)
  }
```



# The Correctness of the Recursive Factorial Method

- Induction on  $n$  can prove that the method `fact` returns the values

$$\text{fact}(0) = 0! = 1$$

$$\text{fact}(n) = n! = n * (n - 1) * (n - 2) * \dots * 1 \quad \text{if } n > 0$$

# The Cost of Towers of Hanoi

- Solution to the Towers of Hanoi problem

```
solveTowers(count, source, destination, spare)
  if (count is 1) {
    Move a disk directly from source to destination
  }
  else {
    solveTowers(count-1, source, spare, destination)
    solveTowers(1, source, destination, spare)
    solveTowers(count-1, spare, destination, source)
  }
```

# The Cost of Towers of Hanoi

- Question
  - If you begin with  $N$  disks, how many moves does `solveTowers` make to solve the problem?
- Let
  - `moves (N)` be the number of moves made starting with  $N$  disks
- When  $N = 1$ 
  - `moves (1) = 1`

# The Cost of Towers of Hanoi

- When  $N > 1$   
$$\text{moves}(N) = \text{moves}(N - 1) + \text{moves}(1) + \text{moves}(N - 1)$$
- Recurrence relation for the number of moves that `solveTowers` requires for  $N$  disks  
$$\text{moves}(1) = 1$$
$$\text{moves}(N) = 2 * \text{moves}(N - 1) + 1 \quad \text{if } N > 1$$

# The Cost of Towers of Hanoi

- A closed-form formula for the number of moves that `solveTowers` requires for  $N$  disks

$$\text{moves}(N) = 2^N - 1, \text{ for all } N \geq 1$$

- Induction on  $N$  can provide the proof that  $\text{moves}(N) = 2^N - 1$

# Summary

- Backtracking is a solution strategy that involves both recursion and a sequence of guesses that ultimately lead to a solution
- A grammar is a device for defining a language
  - A language is a set of strings of symbols
  - A recognition algorithm for a language can often be based directly on the grammar of the language
  - Grammars are frequently recursive

# Summary

- Different languages of algebraic expressions have their relative advantages and disadvantages
  - Prefix expressions
  - Postfix expressions
  - Infix expressions
- A close relationship exists between mathematical induction and recursion
  - Induction can be used to prove properties about a recursive algorithm